

Opening Keynote

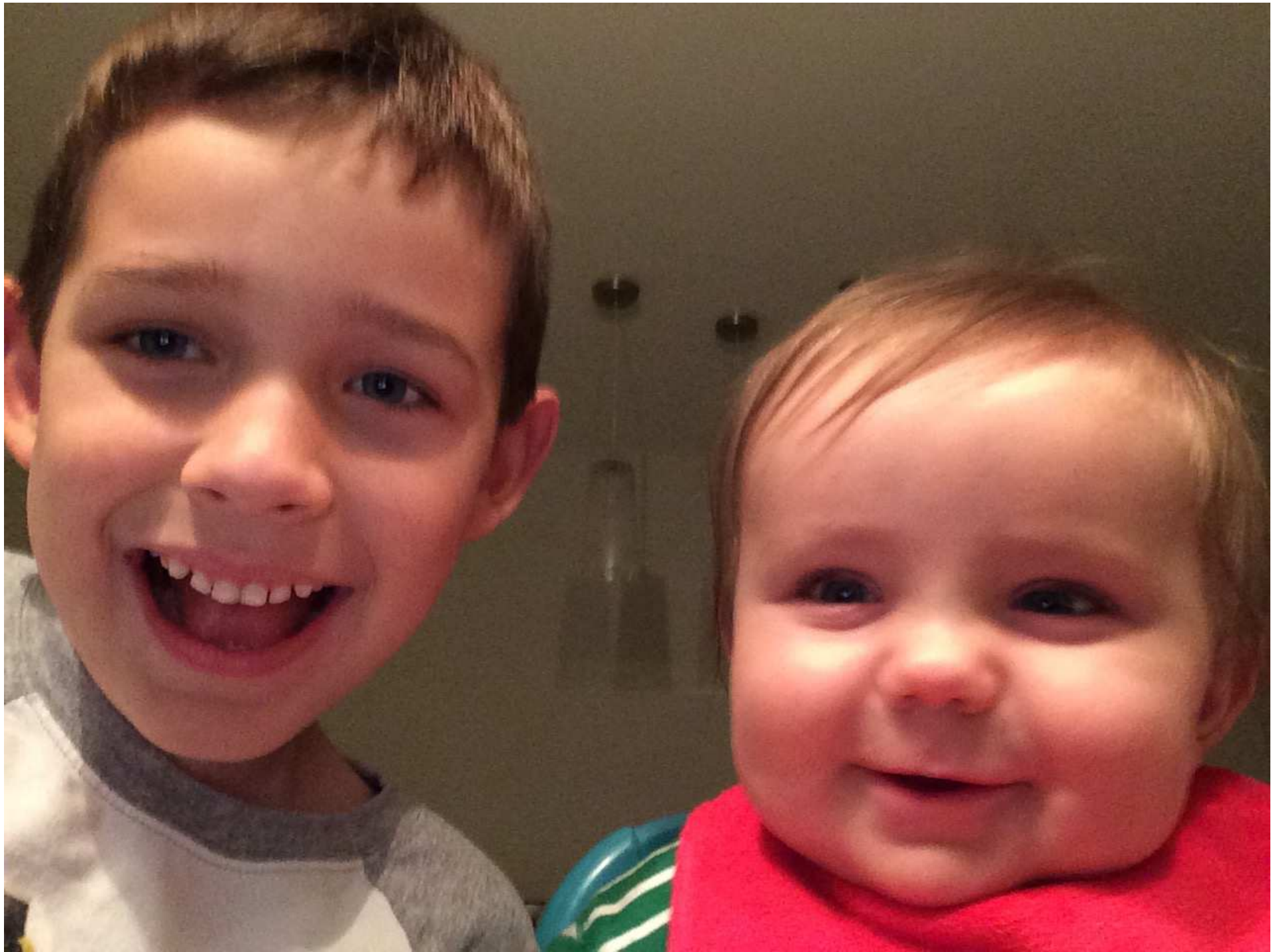
Prepared for DConf 2016

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Welcome to DConf 2016!

Sup



First Five Minutes

- First five minutes become the next five years
- Out of the box experience
 - Website
 - Tutorials
 - Documentation
 - Libraries
- (Self-)Curating dub

Scaling Up

- Many thanks for a great year!
 - Thanks for being here!
 - The Foundation is up!
- Roles: build czar, sysadmin, webmaster, social media, conference organizer...
- Enhance the “point of contact” approach

- Please get me fired!
- I’m in charge of too many things

Raising the Bar on Contributions

- The weird thing about “okay work”
- Reasons to NOT pull something:
 - “Porque no?”
 - Respected contributor
 - “That’s a lot of work”
 - One-liners and many-liners
 - Renames
 - Refactorings showing no clear improvement
 - Churn, illusion of progress

Raising the Bar on Contributions

- Reasons to pull something:
 - Adds value
 - Refactors for a net benefit
 - Fixes a bug “positively”
 - Makes code simpler
 - Makes code faster
 - Automates
 - Improves documentation

- talent × time = win

Resource Management

Take the Garbage Out

Commit to making D
great *with and without*
a GC

Issues with RC

- Overall: RC a qualified success story in computing
- Make it work with `immutable`
- Make it safe
- Make it work with `classes` and value types
- Make it work lazily (COW)
 - Copies should not have arbitrary cost

Basics

- The C++ way: library smart pointers
 - + Apply to any type
 - Apply to any type (unsafe)
 - Needs mutable
- Envisioned:
 - In-language support + library hooks
 - Attribute @rc
 - `opIncRef(uint)`, `opDecRef(uint)`

Safety

- Main issue: unaccounted **refs**

```
struct RC { int a; ... }  
void fun(ref RC obj, ref int x) {  
    obj = RC();  
    ... use x ...  
}  
void gun() {  
    RC obj = ...;  
    fun(obj, obj.f);  
}
```

- Variant: obj is global
- Attack: Insert extra incs/decs for **ref** params
- Fuse successive incs/decs where possible
- Elide increments where possible

immutable

- Apparent contradiction!
 - **immutable**: “I’ll never change, honey”
 - RC: surreptitious change in metadata
- We considered revoking **immutable** rights for RC objects
- If only we had a means to:
 - allocate metadata along with each object
 - type metadata independently
 - access metadata in $O(1)$...

“Why do you completely discard external reference counting approach (i.e. storing refcount in GC/allocator internal data structures bound to allocated memory blocks)?”

– Dicebot

“Never ascribe to malice that
which is adequately explained by
incompetence.”

– Robert J. Hanlon

AffixAllocator!... AffixAllocator!...

- Part of `std.experimental_allocator` from day one
- `AffixAllocator!` (`GCAAllocator`, `uint`):
 - fronts each allocation with an extra `uint`
 - ...that's independently typed
 - Accessible in $O(1)$!

- Use this allocator for creating RC objects

```
import bigo;
```

“Big O” Notation

- Compact characterization of algorithms
 - Growing importance in recent years
 - Usually confined to documentation
 - Pen and paper suffices for non-generic code
-
- Better: make it generic and a discoverable part of the API

Scaling Naming Conventions

- Nomenclature approaches don't scale
- `removeLinTime`, `removeLogTime`,
`removeConstTime`
- Hierarchy of speeds: faster functions
subsume slower ones
- Sometimes $O(\cdot)$ depends on 2+ parameters
- Helpless with HOFs
- Doesn't scale to large APIs

- We want to automate this

Related Work

- Java: initially complexity-oblivious APIs
 - Later: RandomAccess “marker” interface
- STL carefully specifies complexity
 - Archetypal examples: `push_front`, `push_back`, `operator[]`
 - Syntax complexity follows algo complexity
 - Undecided on “best effort” vs. “present or not”, e.g. `distance`

Loosely Related Work

- $O(\cdot)$ analysis part of typechecking (Marion 2011)
- Automated Higher-Order Complexity Analysis (Benzinger 2004)
- Monotonic State (Pilkiewicz et al 2011)

Here

- User:
 - Introduces annotations
 - Defines composition
- Framework:
 - Provides algebra
 - Propagates attributes
 - Calculates composition
 - Notably for higher-order functions
 - Makes result available by static introspection

Synopsis

// Generic doubly-linked list of E

```
struct DoublyLinkedList(E) {
```

```
    ...
```

```
    // Complexity is  $O(1)$ 
```

```
    void insertFront(E x) @O(1);
```

```
}
```

// Generic contiguous array of E

```
struct Array(E) {
```

```
    ...
```

```
    // Complexity is  $O(n)$  in the first argument (this)
```

```
    void insertFront(E x) @O("n");
```

```
}
```


Synopsis

```
// Complexity of insertFrontMany is the complexity of  
// C.insertFront multiplied by the size of the second  
// argument.  
void insertFrontMany(C, E)(ref C container, E[] items)  
@(complexity!(C.insertFront) * O("n2")) {  
    foreach (item; items) {  
        c.insertFront(item);  
    }  
}
```

Introspection

```
static assert(  
    complexity!(insertFrontMany!MyC) <=  
        O("n2") * log(O("n1")),  
    "Too high complexity for insertFrontMany.");
```

Conventions

- Unannotated functions are considered $O(1)$
- "nk" for the k th parameter
- **this** is the first parameter
- "n" if only one parameter of interest

Algebra

- Credit: Timon Gehr

$$C \triangleq O \left(\sum_i \prod_j v_{ij}^{p_{ij}} \log^{l_{ij}} v_{ij} \right)$$

- p_{ij}, l_{ij} positive, $p_{ij} + l_{ij} > 0$
- $\log n, n, n \log n, \sqrt{n_1} + n_2 \log n_2, n_2 \log n_1 \dots$

Normal Form & Partial Order

- Normal form for terms: most compact form
 - $A \leq A'$ immediate (compare vars and powers)
 - $T \leq T'$ iff for each atom A in T there's an atom A' in T' with $A \leq A'$
 - $C \leq C'$ iff for each term T in C there's a term T' in C' with $T \leq T'$
-
- Normal form for complexities: no terms are ordered by \leq

Operations on Complexities

- Comparison for equality and \leq
- Addition (add, then normalize)
- Multiplication (multiply, then normalize)
- Normalization keeps only the fastest-growing terms: $O(n + \sqrt{m} + m \log n) + O(m^2 + \log n)$ is $O(n + m^2 + m \log n)$

- log just a bit trickier (can't express $\log(n_1 + n_2)$)

log for Complexities

- Rely on a simple approximation

$$\log \left(\sum_i \prod_j v_{ij}^{p_{ij}} \log^{l_{ij}} v_{ij} \right) \triangleq \sum_i \sum_{j, p_{ij} > 0} \log v_{ij} \quad (1)$$

- $\log \log$ very slow growing, ignore
- $\log(a + b) \leq \log(ab)$

Implementation

- Operator overloading (`==`, `<=`, `+`, `*`)
- Pivotal use of compile-time evaluation
 - Perfect match with attribute expressions
- Run-time computation automatically available
- Sweet spot between convenience and complexity (sic)

- Coming soon!

One Last Thing