# Using Floating Point Without Losing Your Sanity 

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## Sanity Checks

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- A Crisis of Confidence



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- A Crisis of Confidence
- The specialists, "Numerical Analysts", are rare -- yet ordinary programmers need to use floating point
- It's more fun if you view it as magic


## A Child's Magic Trick

Think of a number ...
Double it
Add 8
Halve it
Take away the number you first thought of And your answer is ...

## An Adult's Magic Trick

- Think of a floating point number...

```
float magic ( float x )
{
return \(x+35-x\);
\}
```

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- "Catastrophic Cancellation"


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We're pulling a trillion rabbits out of a 32-bit hat

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- $35+(1000000000-1000000000)==35$
- Why do we use such a grotesque, fraudulent type?


## Floating point is a success story

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- Total GPU power exceeds CPU power
- Despite being a horrendous approximation, 64 bit floating point is "good enough"


## Two worlds

- The Mathematician's World

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- Sometimes we try too hard to stay in the Mathematician's World


## 3 Misconceptions

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- BELIEF: Floating point is weird
- REALITY: Most real-world measurements are similar


## Floats are just ints with a scale

struct float \{<br>bool sign; int mantissa; int exponent;<br>\}

- mantissa * $2^{\wedge \wedge}$ exponent


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## Floats are just ints with a scale

struct float \{
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- mantissa * 2 ^^ exponent
- If exponent is 0 , it really is an integer
- Most important property is the precision: the number of bits in the mantissa.
- In D, float.mant_dig gives the precision


## The Precision Budget

The larger the precision, the more extravagent you can be

| float | 22 bits |
| :--- | :--- |
| double | 54 bits |
| real | 64 bits |
| quadruple | 112 bits |


| Operation | Cost |
| :--- | :--- |
| Multiplication | 1 bit |
| Division | 1 bit |
| Addition | Many |
| Take away the number you first thought <br> of | Bankrupt |



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## Floating Point Exceptions

- The hardware can generate hardware traps when funny values are produced. Most programs should enable the severe traps inside main()

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- Unfortunately there is no way to detect Catastrophic Cancellation
"Don't use =="
- Why not? Because it destroys the illusion



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- Some implementation details are hidden
- == is still useful for low-level code and unittests.


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- In Physics, there is no "exact equality" either
- Always need to specify the precision


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- But sometimes we have Orwellian experiences...



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- assert( $1.30==1.30 f$ ); // OK!!


## Sociomantic's Nine Trillion Dollar Bug

- Losing your sanity, \#1

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if ( price < 0 ) { error(); }
if ( price ) {
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- DMD Issue \#13489 - never do an implicit cast from float to bool unless you can guarantee it is not NaN .


## Generic Programming

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## Generic Programming

- Mathematically, reals are an extension of integers
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- Replace 'int' with 'double' and everything will compile
- Test cases will still work
- So let's make our code work with any numeric type!



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- The VALUES are a superset of int
- The SEMANTICS are not
- For generic code we need common semantics



## Losing Your Sanity, \#2

- A simple foreach range
int doTen ( T ) ( T from )
\{
int howmany $=0$;
foreach ( x ; from .. from + 10)
++howmany;
return howmany;
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- doTen!float( 500 ) == 10
- doTen!float( 16777242 ) == 9
- doTen!float( 18000000 ) does not terminate


## Increment (or not)

For integers, ++x; --x; is a no-op
For floats it's more fun

| $\mathbf{x}$ | After ++x; --x; |
| :--- | :--- |
| 31837 | 31837 |
| $1.25 \mathrm{e}-6$ | $1.20 \mathrm{e}-6$ |
| $-1 \mathrm{e}-20$ | 0 |
| 16777250 | 16777252 |

- If you use ++ on a float, someone will go insane.


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- All uses of isNumeric() are trivial, except two
- std.complex just casts integers to floating point
- std.random.dice() is incorrect for pathological cases
- There are probably no mathematical algorithms that work for both integers and floating point


## "More Precision Is Always Better"

- More precision improves the illusion.
double magic ( double $x$ )
\{
return $x+35-x$;
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double magic ( double $x$ )
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- magic(1000000000) $==35$
- magic(5e17) == 64
- Corner cases move but don't disappear


## Rounding Modes

| Rounding <br> Mode | $\mathbf{2 . 5}$ | $\mathbf{- 5 . 5}$ |
| :--- | :--- | :--- |
| Round to Near- <br> est | 2 | -6 |
| Round Up | 3 | -5 |
| Round Down | 2 | -6 |
| Round To Zero | 2 | -5 |



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- Double rounding is a problem.
- 3.49 rounds down to 3
- 3.49 rounds up to 3.5 , which rounds up to 4


## Secret Precision

- Extra hidden precision can happen when:

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- Extra hidden precision can happen when:
- The $x 87$ FPU is used on $x 86$ machines


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- real only has 64 bits. We'll round twice.
- One in 1024 calculations has an out-by-1 error


## In Practice

Most library code splits the possible input values into smaller ranges, and then performs a different calculation for each range

## Root finding

Given a function double $\mathrm{f}($ double x ), $\mathrm{f}(\mathrm{x} 0)>0, f(\mathrm{x} 1)<0$
find the point where $f(x)==0$

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- With 80 -bit reals, worst case is $>16000$ calls


## The Binary Chop That Isn't

auto midpoint $=(x 0+x 1) / 2$;

- Let $x 0=1 \mathrm{e} 100, \mathrm{x} 1=1 \mathrm{e}-100$, and ultimate solution is $2 \mathrm{e}-100$


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- We get to 2e-100 after 600 iterations



## Binary Chop For Real

- Midpoint in implementation space
ulong x0_raw = reinterpret!ulong(x0);
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auto midpoint $=$ reinterpret!double( x0_raw + x1_raw ) / 2;



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- Again let $x 0==1 e 100, x 1=1 e-100$, and solution is $2 e-100$
- Midpoints are 5e0, 2.5e-50, 1.2e-75, 6e-88, 3e-94 ...
- We reach $2 \mathrm{e}-100$ after 9 iterations


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- TOMS 748 has a similar problem with linear interpolation
- Fixing that improves the average case as well.
- Available in std.numeric.findRoot



## Moral

Even when floating point code compiles, and gives the mathematically correct answer, it can still be algorithmically wrong

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- float requires great care. Prefer double or real.
- Use == only when you want to expose implementation details
- Generic numeric code is almost certainly wrong in horrible, subtle ways
- D is (mostly) a pleasant language for floating point.


## Questions?


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