Unified Concurrent Runtime for D

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Unify what?

In essence it's all about fibers
...and threads that schedule them

Fiber is a cooperative user-mode thread

D's fibers are scheduled *manually*

Hence we have a multitude of custom schedulers

Vibe.d is a popular one
```c
struct Node {
    int payload;
    Node* left, right;
}

void main(){
    int current;
    auto tree = new Node(3, new Node(2, new Node(1)), new Node(4));
    void postorder(Node* n) {
        if (n.left) postorder(n.left);
        current = n.payload;
        Fiber.yield();
        if (n.right) postorder(n.right);
    }
    Fiber f = new Fiber(() => postorder(tree));
    for (;;) {
        f.call();
        if (f.state == Fiber.State.TERM) break;
        writeln(current);
    }
}
```
Fibers and I/O

Now with I/O it gets tricky - the thread would be blocked

```c
// Sketch of TCP echo server
import std.socket;
void echo(Socket sock, ubyte[] buffer)
{
    ssize_t read, write;
    for(;;) {
        do {
            read = sock.read(buffer[]);
            if (sock.wouldHaveBlocked) Fiber.yield();
            else if (read <= 0) break;
        } while(read < 0);
        do {
            write = sock.write(buffer[0..read]);
            if (sock.wouldHaveBlocked) Fiber.yield();
            else if (read <= 0) break;
        } while(write < 0);
    }
}
```

Suboptimal w/o eventloop + somebody needs to 'call' fibers
Fibers and I/O

Vibe.d chooses to provide its own I/O API

```d
import vibe.vibe;

void main()
{
    listenTCP(7, (conn) { conn.write(conn); });
    runApplication();
}
```

Nice and clean

Or is it?
Caveats

Using anything not Fiber-aware will block thread and ruin performance

There are zero C/C++/Rust libraries that are made to run on top of Vibe.d

std.net.curl, std.socket, std.process do not work with Fibers
std.concurrency was extended to work with custom scheduler

To Fiberize or not, that is the question
Project Photon

Transparent fiber scheduler that works with *most* blocking C/C++/etc code

Multi-threaded with great performance out of the box

Look beyond I/O and 3rd party libraries: Events, Channels/Queues
Photon in action

Parallel download of files with std.net.curl, Photon edition

```pascal
import std.net.curl, std.stdio, std.datetime.stopwatch;
import photon;

immutable urls = [ "https://...", ... ];

void main()
{
    startloop(); // init Photon eventloop (to be done lazily soon)
    void spawnDownload(string url, string file) {
        spawn(() => download(url, file));
    }
    StopWatch sw;
    sw.start();
    foreach(url; urls) spawnDownload(url, url.split('/').back);
    runFibers(); // start schedulers and run fibers
    sw.stop();
    writefln("Done: %s ms", sw.peek.total!"msecs");
}
```
## Photon in action

Parallel download of files with std.net.curl, fiber edition

<table>
<thead>
<tr>
<th></th>
<th>12 files (~1Mb file)</th>
<th>3 files (~200Kb each)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sequential:</td>
<td>16 sec</td>
<td>4.4 sec</td>
</tr>
<tr>
<td>Threads:</td>
<td>5.3 sec</td>
<td>0.46 sec</td>
</tr>
<tr>
<td>Fibers:</td>
<td>5.9 sec</td>
<td>0.52 sec</td>
</tr>
</tbody>
</table>

The test is super noisy (easily ~30%) on the internet.

We are in the ballpark of dedicated threads...even when running photon on a single core.
Transparent scheduler

Fiber.yield needs to happen in some 3rd party library

Most code calls libc syscall wrapper!

That is precisely the point of our "integration" ... on Posix-like OS
Shadowing libc syscalls

Put in a shared library linked \textit{before} the libc

\begin{verbatim}
extern(C) ssize_t read(int fd, void *buf, size_t count) noexcept
{
    return universalSyscall!(SYS_READ, "READ", SyscallKind.read,
                            Fcntl.explicit, EWOULDBLOCK)
                 (fd, cast(size_t)buf, count);
}

... // 6 or so more

ssize_t universalSyscall(size_t ident, string name, SyscallKind kind,
                          Fcntl fcntlStyle, ssize_t ERR, T...)
     (int fd, T args) noexcept {
    if (currentFiber is null) {
        logf("%s PASSTHROUGH FD=%s", name, fd);
        return syscall(ident, fd, args).withErrorno;
    }
    else { ... MAGIC IS HERE }
}
\end{verbatim}
Scheduler (Linux ver.)

Thread per "Core" (SMT etc.) with affinity mask

All I/O state is static and pre-mapped memory

"Cowboy" lock-free scheduler(s) <-> event loop interface

Fibers are pinned to thread based on load via "power of 2 random choices"

Some things are "Research code" quality
Scheduler in picture (v2)

Run Queue

Scheduler Core 1

Run Queue

Scheduler Core 2

Run Queue

Scheduler Core 3

Event loop

No more fibers?
Scheduler in picture (v3?)

- Run Queue
  - Scheduler Core 1
  - Scheduler Core 2
  - Scheduler Core 3

- EV Thread Any Core
  - Event loop
Descriptor state

- NONBLOCKING
  - close
  - socket-like

- NOT INITED
  - close
  - first syscall

- THREAD POOL
  - close
  - file-like

- INITIALIZING
  - first syscall
  - socket-like
Descriptor read state

- **EMPTY**
  - Schedule waiters

- **READY**
  - Schedule waiters

- **UNCERTAIN**
  - May become uncertain if read full buffer

- **READING**

**SCHEDULER**

**EVENT-LOOP**
Implementing `poll`

- FD #1
- FD #2
- FD #3

Event ➔ Fiber 1 ➔ Fiber 2

Steal the fiber
Implementing `poll` (2)
What about Windows?

Same technique would work...

but there is no "small" libc

WinAPI is 10,000+ API entry points

Yet there is User Mode Scheduling subsystem

Which is kernel-assisted support for user-mode threads (of sorts...)
UMS in 5 seconds

Under User-Mode we get to create real threads:

1. *Never* scheduled by the NT *kernel scheduler*

2. To run them we pick a set of normal dedicated (per core typically) *scheduler threads and run them explicitly*

Better then a fiber in that:

1. An OS will automatically switch to scheduler should the current thread about to block

2. Any *blocking event such as pagefault*

3. *Each thread has 1:1 user:kernel mapping, TLS and other goodies*
UMS in pictures

Kernel

Userland

KT0

KT1

KT2

UT0

UT1

UT2

UMS scheduler thread

Without switching the KT-s
UMS in pictures

Eventually we do a syscall and NT restore sanity
UMS in pictures

We switch away to UT2
UMS in pictures

KT0  KT2  KT1

UT2
UMS scheduler thread

UT0

UT1

KT1 is completed and UT1 is put in completion queue where our scheduler grabs it.
It's not all roses

On subject, there is

1 interview with Windows Kernel Architect

2 examples with comments like "meh, doesn't improve performance for simple test"

There was nice video at the BUILD conference on Windows 8, can't find it anymore

The real kicker:
Not my code nor any of 2 examples actually run on VirtualBox VM
"A few months ago, as part of looking through the changes in Windows 10 Anniversary Update for the Windows Internals 7th Edition book, I noticed that the kernel began enforcing usage of the CR4[FSGSBASE] feature ... in order to allow usage of User Mode Scheduling (UMS)."

Freaking great find, Alex!

VirtualBox doesn't propagate that CPU flag

For future poor souls seeking to enable UMS:
Closing notes on UMS

It is still more efficient to just use normal threads if we do not do async I/O with completion ports.

Photon overrides parts of WinSock like on Linux to use Overlapped I/O with Completion Ports.

There is a nice option to use RIO sockets, should be super fast but trickier to get right.

UMS thread is ~100x slower to spawn then Fiber.

The whole thing is x64 only and Ivy Bridge or later CPU.
Fast HTTP is not the main point of Photon

Using thin-wrapped Node.js (Nginx) HTTP parser

Using plain std.socket to prove the point

Add in a few sensible performance optimisations

Peper it by a couple of speedhacks

(competion does it as well, so why not)
Case study: HTTP server

HTTP server in Photon runtime

```java
void server() {
    Socket server = new TcpSocket();
    server.setOption(SocketOptionLevel.SOCKET, SocketOption.REUSEADDR, true);
    server.bind(new InternetAddress("0.0.0.0", 8080));
    server.listen(1000);
    void processClient(Socket client) {
        spawn(() => server_worker(client));
    }
    while(true) {
        try {
            Socket client = server.accept();
            processClient(client);
        }
        catch(Exception e) {
            writef("Failure to accept \%s", e);
        }
    }
}
void main() {
    startloop();
    spawn(() => server());
    runFibers();
}
```
Case study: HTTP server

HTTP server in plain D threads

```d
void server() {
    Socket server = new TcpSocket();
    server.setOption(SocketOptionLevel.SOCKET, SocketOption.REUSEADDR, true);
    server.bind(new InternetAddress("0.0.0.0", 8080));
    server.listen(1000);
    void processClient(Socket client) {
        new Thread(() => server_worker(client)).start();
    }
    while(true) {
        try {
            Socket client = server.accept();
            processClient(client);
        } catch(Exception e) {
            writeln("Failure to accept %s", e);
        }
    }
}
void main() {
    new Thread(() => server()).start();
}
```
Case study: HTTP server

Common part - bare bones HTTP processor

```java
import utils.http_server;

class HelloWorldProcessor : HttpProcessor {
    HttpHeaders[] headers = [HttpHeader("Content-Type", "text/plain; charset=utf-8"

    this(Socket sock) { super(sock); }

    override void onComplete(HttpRequest req) {
        respondWith("Hello, world!", 200, headers);
    }
}
```

Time to bench threads vs fibers
And a few of top guns from TechEmpower benchmark
And I'm taking DMD to this gunfight
HTTP server benchmark

Photon vs Threads
HTTP hello

Concurrency

RPS Photon
RPS threaded
Latency 99% Photon
Latency 99% Threaded
HTTP server benchmark

Graph showing the performance of different HTTP servers under varying concurrency levels.

- Photon
- D threaded
- Fast HTTP Go
- Undertow (Java)

Y-axis: Requests per second (RPS)
X-axis: Concurrency
HTTP server benchmark

All things latency
HTTP hello

Concurrency

Photon
D threaded
Fast HTTP Go
Undertow (Java)
HTTP server benchmark

All things RSS(mb)
HTTP hello

Concurrency

Concurrency

Photon
D threaded
Fast HTTP Go
Undertow (Java)
Beyond I/O

Channels or queues got popular with Go

OutputRange!T

Event object

InputRange!T

Just label the obvious. Works for MPSC queue
Event, channel, source

```cpp
interface Event {
    bool ready();
    void reset(); // typically called by await in Event-loop
}

size_t await(R)(R events)
if(isEvent!ElementType!R)) {
... event loop / scheduler magic
}
```

Channel is an OutputRange

Source is an InputRange
Event-aware algorithms

```plaintext
struct Map(R) {
    R _range;

    bool ready() { _range.ready: }

    auto event(){ return _range.event; }
}

struct Zip(RS...) {
    RS _ranges;

    bool ready() {
        return _ranges.reduce!(true)((a,b) => a && b.ready);
    }

    auto event(){
        return _ranges.map!(x => x.event);
    }
}
```
Bright future

Most of the ground work is done on:
- basic I/O, primitive scheduling of Fibers
- transparent integration of C libraries shown to work

Time to capitalize and expand:
- define composable multiplexing patterns
- "Executors" with Photon's current scheduler as an option
- Vibe.d on Photon (?)
- expand to cover more syscall surface
Help wanted!

Current platforms - Windows x64 and Linux x64

Even if you are not thrilled to hack on Windows User Mode Something...

Lots of simple work to test things out!

32-bit Linux, FreeBSD and MacOS

BetterC Fibers would be an awesome addition

https://github.com/DmitryOlshansky/photont